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## Background

**Safe Region:**  $\mathcal{C} = \{x: h(x) \geq 0\} \subseteq \mathcal{X}$

**Positive invariance:**  $x(t) \in \mathcal{C} \subseteq \mathcal{X}$  for all  $t \geq 0$ , if  $x(0) \in \mathcal{C}$ .

**Safety:** positive invariance of a given safety region  $\mathcal{C}$ .

**Dynamics:**  $\dot{x}(t) = f(x(t)) + g(x(t))u(t)$  (1)

**Control Barrier Function (CBF)**  $b(x)$  is a smooth function that evaluates the ‘safety’ of the system. Let  $\mathcal{D} = \{x: b(x) \geq 0\} \subseteq \mathcal{C}$ .

Function  $b(x)$  is a CBF if there exist  $u$  and class- $\kappa$  function  $\alpha$ , s.t.

$$\frac{\partial b}{\partial x}(f(x(t)) + g(x(t))u(t)) \geq -\alpha(b(x(t)))$$
 (2)

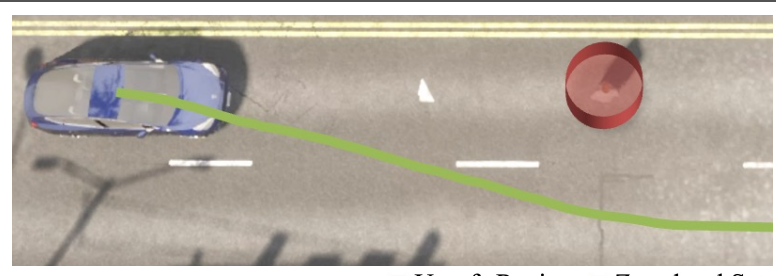
However,  $u$  does not exist if  $\frac{\partial b}{\partial x}g(x(t)) = 0$  and  $\frac{\partial b}{\partial x}f(x(t)) < 0$ .

**Exact Verification of ReLU NCBFs [1]**

**Nagumo's Theorem** The closed set  $\mathcal{D}$  is positive invariant iff, whenever boundary states  $x \in \partial\mathcal{D}$ ,  $u \in \mathcal{U}$  satisfies

$$\frac{\partial b}{\partial x}(f(x) + g(x)u) \geq 0.$$

**Intuition: Piece-wise Linearity of ReLU**



## Problem Studied

Given a system (1), synthesize a ReLU NCBF  $b_\theta(x)$  s.t.

- Feasible: for all  $x \in \mathcal{D}$ , there exist  $u$  satisfying (2)
- Correct:  $\mathcal{D} \subseteq \mathcal{C}$  (safe region)

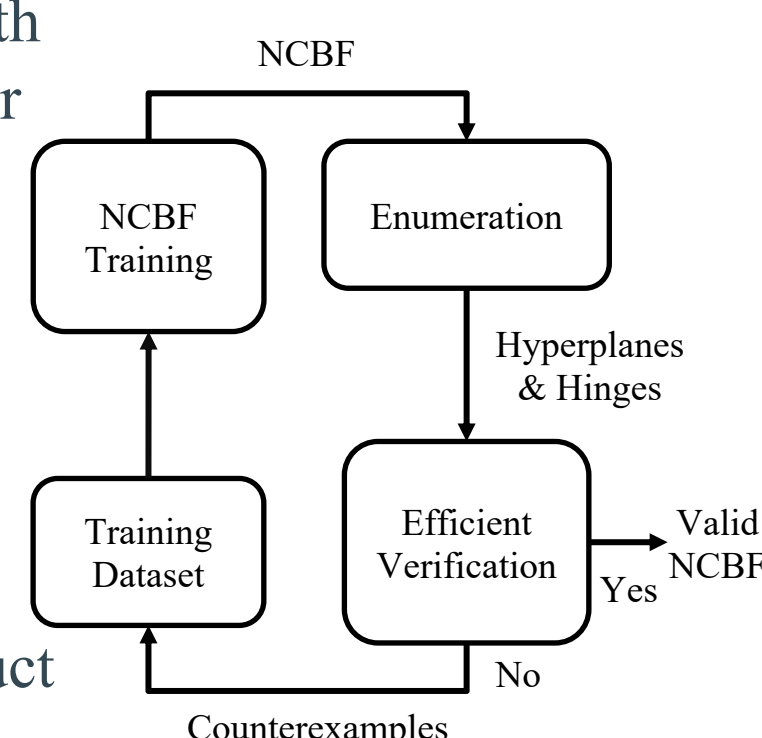
**Challenge 1:** Scalability for high-dimensional systems and deep NNs.

**Challenge 2:** Synthesize NCBFs satisfying conditions

- Loss may not converge to zero
- Hard to obtain a formal verification

## Contributions

1. Propose a framework for Synthesis with Efficient Exact Verification (SEEV) for ReLU NCBFs.
2. Develop a training procedure that reduces the number of segments that must be verified.
3. Construct hierarchical verification that efficiently enumerate hyperplanes, exploit sufficient conditions and conduct exact verification.



## Synthesis with Efficient Exact Verification (SEEV)

Goal: Synthesize NCBFs to be feasible and correct.

Training dataset  $\mathcal{T}$ : initialized by uniform sampling over  $\mathcal{X}$ .

$$\min_{\theta} \lambda_B \mathcal{L}_B(\mathcal{T}) + \lambda_f \mathcal{L}_f(\mathcal{T}) + \lambda_c \mathcal{L}_c(\mathcal{T})$$

$\mathcal{L}_c$  correctness regularizer enforces the correctness of the NCBF.

$\mathcal{L}_f$  feasibility regularizer  $\mathcal{L}_f = \|u - \pi_{nom}(x)\|_2^2 + r$  where  $r$  is the slack variable for the safety filter [2]

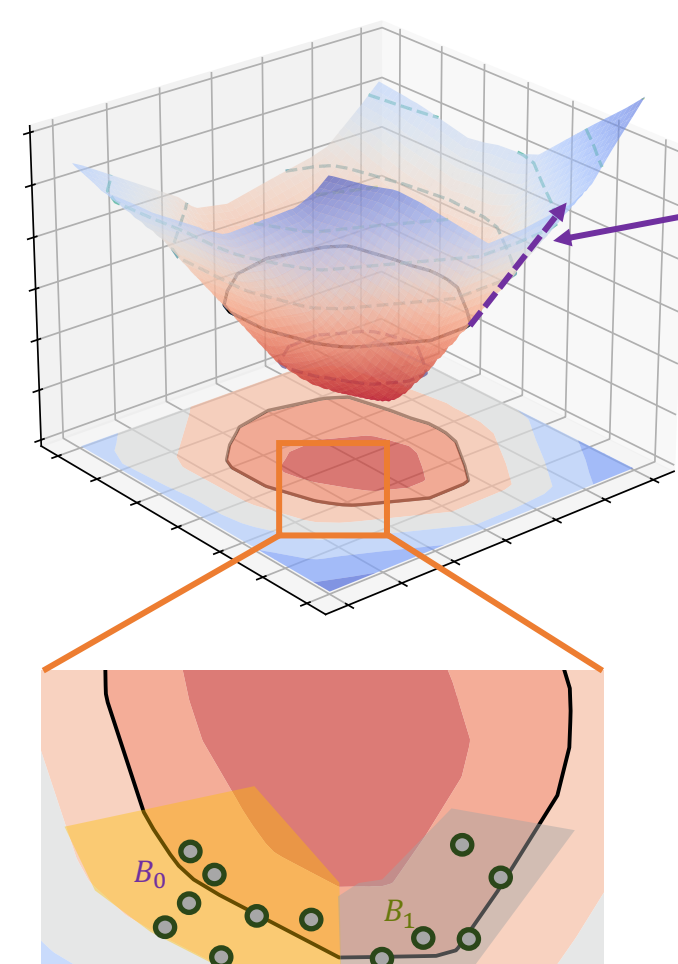
$$\min_u \|u - \pi_{nom}(x)\|_2^2$$

s.t.  $\mathcal{W}(S_i)^T(f(\hat{x}) + g(\hat{x})u) + r \geq 0$

$\mathcal{L}_B$  Boundary activation regularizer limits the number of hyperplanes & hinges by penalizing the dissimilarity

**Step 1** Identify boundary samples; **Step 2** Clustering; **Step 3** Penalize dissimilarity

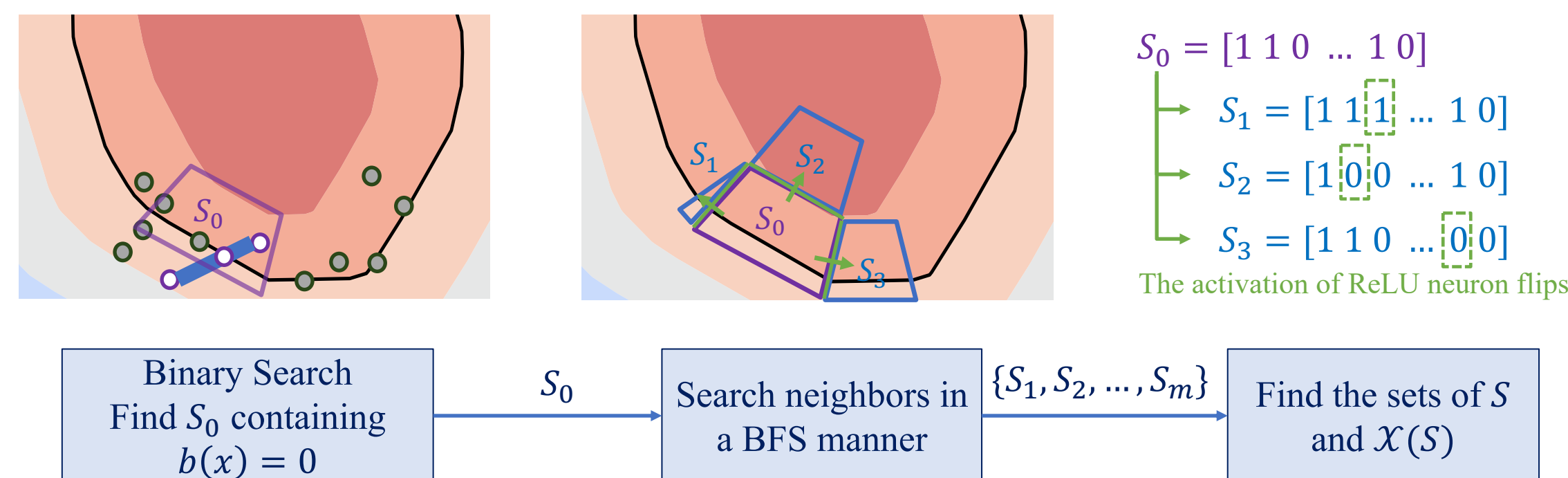
$$\mathcal{L}_B = \frac{1}{N_B} \sum_{B_i \in \mathcal{B}} \frac{1}{|\mathcal{J}_{B_i}|^2} \sum_{\hat{x}_i, \hat{x}_j \in \mathcal{J}_{B_i}} \|\phi_{\sigma_k}(x_i) - \phi_{\sigma_k}(x_j)\|_2^2$$



## Efficient Exact Verification for ReLU NBFs

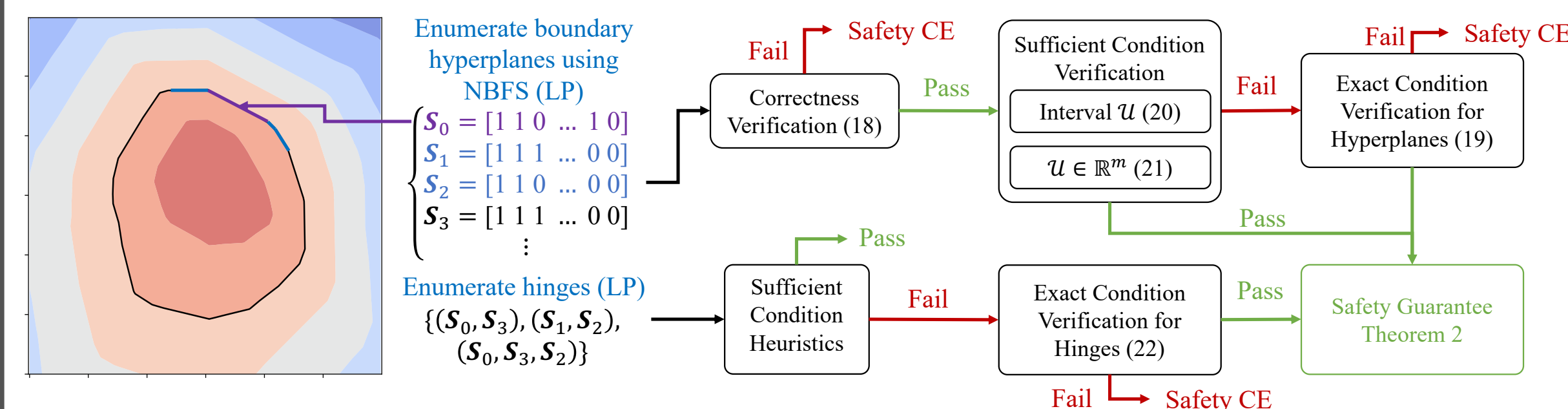
How to compute the  $\partial b/\partial x$  efficiently?

- Derivative of a NCBF can be characterized by activation sets  $\mathbf{S}$ . (Enumerate  $\mathbf{S}$ )
- Derivative of a NCBF in one activation set  $\mathbf{S}$  is linear. (Linear Program)



**Features:**

- Enumeration: (i). Only rely on linear program; (ii). CPU only; (iii). Multi-process enabled
- Verification: (i). Efficient; (ii). Exact verification with SMT solver [3]



## Experiments

We consider the Darboux, Obstacle Avoidance (OA), hi-ord<sub>g</sub> and Spacecraft Rendezvous (SR) problems and compared our approach with SMT-based verifiers.

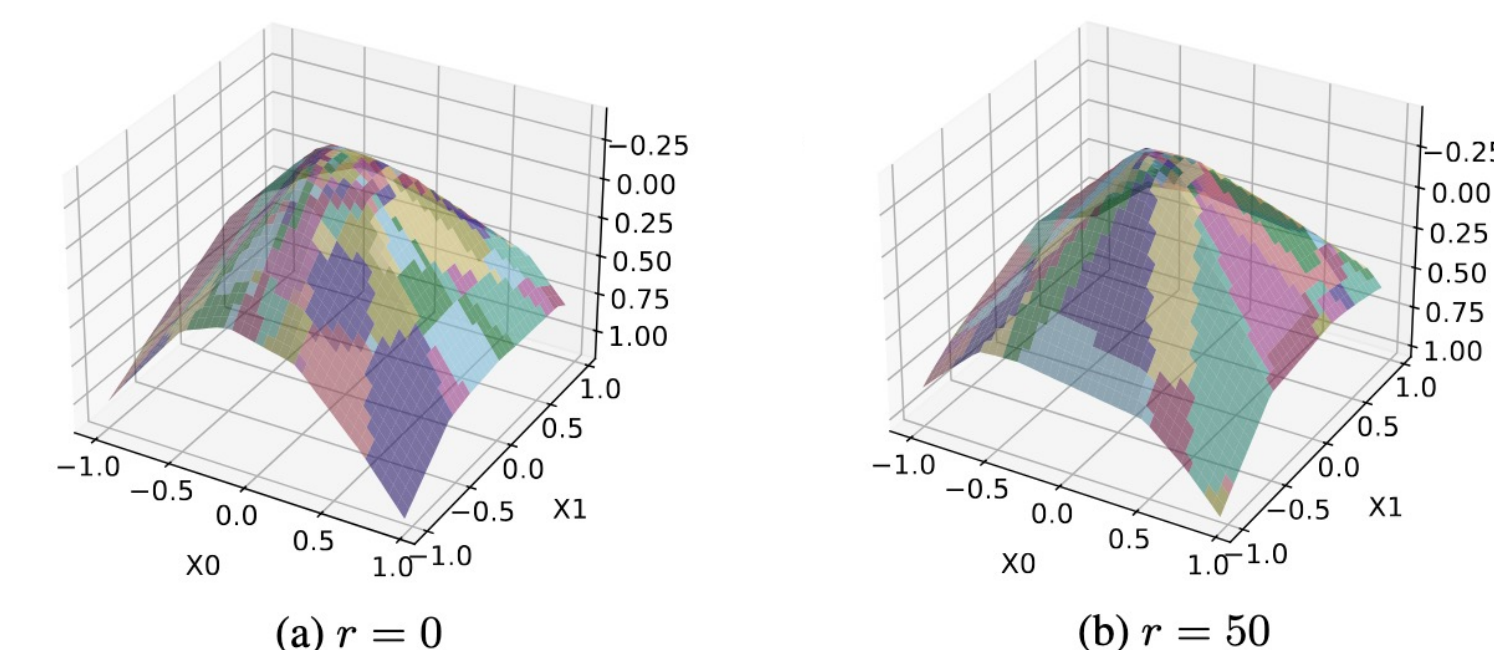


Table1: Comparison of  $N$  the number of boundary hyperplanes and  $C$  coverage of the safe region  $\mathcal{D}$  of NCBF trained with ( $r$ ) and without ( $o$ ) boundary hyperplane regularizer

Case	n	L	M	$N_o$	$C_o$	$N_{r=1}$	$\rho_{r=1}$	$N_{r=10}$	$\rho_{r=10}$	$N_{r=50}$	$\rho_{r=50}$
OA	3	2	8	26	89.46%	25	0.996	23.3	0.994	<b>13.3</b>	1.006
	3	2	16	116	83.74%	119	1.012	111	1.005	<b>98</b>	1.055
	3	4	8	40	91.94%	38	0.988	36	0.993	<b>13</b>	0.937
	3	4	16	156	87.81%	170	0.971	147	1.003	<b>64</b>	1.038
SR	6	2	8	2868	98.58%	2753	1	1559	1	<b>418</b>	1
	6	4	8	6371	98.64%	6218	1	3055	1	<b>627</b>	1
	6	2	16	N/A	N/A	204175	N/A	68783	N/A	<b>13930</b>	N/A

Table2: Comparison of verification run-time of NCBF in seconds.

Case	n	L	M	$N$	$t_h$	$t_g$	SEEV	Baseline [23]	dReal	Z3
Darboux	2	2	256	15	2.5s	0	2.5s	315s	>3h	>3h
	2	2	512	15	3.3s	0	3.3s	631s	>3h	>3h
OA	3	2	16	86	0.41s	0	0.41s	16.0s	>3h	>3h
	3	4	8	15	0.39	0	0.39	16.1s	>3h	>3h
	3	4	16	136	0.65s	0	0.65s	36.7s	>3h	>3h
hi-ord <sub>g</sub>	3	1	128	5778	20.6s	0	20.6s	207s	>3h	>3h
	8	2	8	73	0.54s	0	0.54s	>3h	>3h	>3h
	8	2	16	3048	11.8s	0	11.8s	>3h	>3h	>3h
SR	8	4	16	3984	22.4s	0	22.4s	>3h	>3h	>3h
	6	2	8	2200	7.1s	2.7s	9.8s	179s	UTD	UTD
	6	4	8	4918	45.8s	14.3s	60.1s	298.7s	UTD	UTD

SEEV outperforms the LiRPA-based method proposed in the baseline [23]

SOTA SMT-based Methods are not directly applicable

[1] Zhang, Hongchao, et al. "Exact verification of relu neural control barrier functions." Advances in neural information processing systems 36 (2023): 5685-5705.  
 [2] Dawson, Charles, Sicun Gao, and Chuchu Fan. "Safe control with learned certificates: A survey of neural lyapunov, barrier, and contraction methods for robotics and control." IEEE Transactions on Robotics 39.3 (2023): 1749-1767.  
 [3] Gao, Sicun, Soonho Kong, and Edmund M. Clarke. "dReal: An SMT solver for nonlinear theories over the reals." International conference on automated deduction. Berlin, Heidelberg: Springer Berlin Heidelberg, 2013.

